

- TIM 125/225

- LECTURE #12 (2/13/14)

Derivation of the equations  
for Tailored Aggregation

# 1. Aggregation for Cycle Inventory

3 cases (Multiple Products)

## Case (a) No aggregation

Treated each product separately  
(same as the single product)

Optimal shipment frequency,  $n_i^* = \sqrt{\frac{h_i C_i D_i}{2 S_i^*}} \rightarrow (a)$

$i = 1, 2, \dots, N$  (# of products)

What does  $h$  depend on?

- product life cycle

- size (space)

⋮

## Case (b) Simple aggregation

Each shipment contains all products

$$S^* = \underbrace{S}_{\text{cost of each shipment}} + \sum_{i=1}^N \underbrace{S_i}_{\text{product specific cost}}$$

The optimal  
shipping frequency  
(common to all  
products)

$$n^* = \sqrt{\sum_{i=1}^N \frac{h_i C_i D_i}{2S^*}}$$

↳ 1(b)

### Case (C): TAILORED Aggregation

Used when there are large differences  
in the demands of the various products

e.g.  $D_1 = 10^4$  ;  $D_2 = 10^6$  ;  $D_3 = 10^2$

### Questions

What are the relative shipping  
frequencies for each product?

Idea: When we know the relative frequencies,  
then we can allocate shipping costs  
for each product in a "fair" manner.

## Procedure:

Step 1: Treat each product as independent with respect to shipping (Case a)

Calculate the shipping (ordering) frequency

for product  $i$ ,  $\bar{n}_i = \sqrt{\frac{h_i C_i D_i}{2(S + s_i)}} \rightarrow (2)$

$\uparrow$  fixed       $\uparrow$  variable

$i = 1, 2, \dots, N \triangleq$  total # of products

Determine the product  $k$  for which the shipping frequency is a maximum

i.e.  $\bar{n}_k \triangleq \max(\bar{n}_1, \bar{n}_2, \dots, \bar{n}_N) \rightarrow (3)$

Example if we have 3 products

$$\bar{n}_1 = 30, \quad \bar{n}_2 = 100, \quad \bar{n}_3 = 12$$

then  $\bar{n}_k = \bar{n}_2 = 100 \quad \& \quad k = 2$

Define  $\bar{n} = \bar{n}_k = \max(\bar{n}_1, \bar{n}_2, \dots, \bar{n}_N) \rightarrow (4)$

shipping frequency of  
the most frequently ordered product

Step 2 :

Since each shipment will contain product  $k$ , let's allocate the fixed shipping cost  $S$  to product  $k$

for product  $k$ ,  $S_k^* = S + s_k \rightarrow (5)$

$\uparrow$  fixed       $\uparrow$  variable

for all other products,  $i \neq k$ ,  $S_i^* = s_i \rightarrow (6)$

$(i=1, 2, \dots, N; i \neq k)$

Step 3 :

Recalculate the order frequencies,  $\bar{n}_i$   
for each product based on the  
shipping costs in (5), (6)

Therefore  $\bar{n}_k = \frac{1}{N} \sqrt{\frac{h_k C_k D_k}{2 S_k^*}} \stackrel{(5)}{=} \sqrt{\frac{h_k C_k D_k}{2(S+S_k)}} \stackrel{(2)}{=} \bar{n}_k \stackrel{(4)}{=} \bar{n} \rightarrow (7)$

for all other products ( $i \neq k$ )

$$\bar{n}_i = \frac{1}{N} \sqrt{\frac{h_i C_i D_i}{2 S_i^*}} \stackrel{(6)}{=} \sqrt{\frac{h_i C_i D_i}{2 S_i}} \rightarrow (8)$$

$(i = 1, 2, \dots, N; i \neq k)$

(The products are still being "independently shipped")

step 4: Determine the relative shipping frequency,  $m_i$ , for each product

$\bar{n}_k$   $\leftarrow$   $\bar{m}_i = \left[ \frac{\text{(shipping frequency of the most frequently shipped product)}}{\text{shipping frequency of the } i^{\text{th}} \text{ product}} \right] \stackrel{(7,8)}{=} \frac{\bar{n}_k}{\bar{n}_i} \checkmark$

$$\bar{m}_i \geq 1 \quad (i = k, m_i = 1)$$

Next

$$m_i \triangleq \lceil \bar{m}_i \rceil, \quad i = 1, 2, \dots, N \rightarrow (10)$$

(round-up  
to the nearest  
integer)

$$m_k = \lceil \bar{m}_k \rceil = 1 \quad (\text{most frequently ordered product})$$

In our example,

$$\bar{m}_1 = \frac{100}{30} = 3.3; \quad \bar{m}_2 = \frac{100}{100} = 1; \quad \bar{m}_3 = \frac{100}{12} = 8.3$$

$$m_1 = \lceil \bar{m}_1 \rceil = 4, \quad m_2 = 1, \quad m_3 = \lceil \bar{m}_3 \rceil = 9$$

Step 5: Tailored Aggregation

Allocate shipping costs per order based on the relative shipping frequencies,  $m_i$

$$S^* = S + \frac{S_1}{m_1} + \frac{S_2}{m_2} + \dots + \frac{S_N}{m_N} \rightarrow (11)$$

cost per shipment      ↑ fixed cost

Step 6 : Pose the optimization problem to minimize total costs (material cost + shipping cost + inventory holding cost) as in Case 2 or Case(b) [Simple aggregation]

$$\left[ m_i = \frac{n_k}{n_i} \right] \Rightarrow n_i \stackrel{\text{most frequent}}{=} \frac{n_k}{m_i}, \quad (i=1, 2, \dots, N) \rightarrow (12)$$

(total cost)  $C_T \stackrel{(12)}{=} f(n_k)$

$$\frac{dC_T}{dn_k} = 0 \Rightarrow n_k^* \stackrel{(12)}{=} \sqrt[N]{\frac{\sum_{i=1}^N h_i C_i D_i m_i}{2S^*}} \rightarrow (13)$$

Optimal shipping frequency of the most frequently ordered product

where  $S^*$  is given by (11)

Shipping frequency,  $n_i^*$  for all other products

$$(12) \quad \frac{n_k^*}{m_i} \rightarrow (14) \quad (i=1, 2, \dots, N)$$

Optimal lot size  $(Q_L)_i^*$

$$= \frac{D_i}{n_i^*} \rightarrow (15) \quad (i=1, 2, \dots, N)$$

Use this procedure in  
 HW #5, ~~Prob 4~~  $\Rightarrow$  Ex 10.3 (Third Edition)  
 Problem 3 (EX 10.3 Text)  
 Harley-Davidson  
 done in HW#4 { (1) Separate Products  
 (2) Simple Aggregation  
 (3) Tailored Aggregation

How was eqn (B) for  $n_k^*$  obtained?

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$$C_i' = \begin{array}{c} \text{Transportation} \\ \text{cost} \end{array} + \begin{array}{c} \text{Inv holding} \\ \text{cost} \end{array}$$

$$= (S^*)(n_k) + \sum_{i=1}^N \left(\frac{Q_L}{2}\right)_i h_i C_i$$

$$C_i' = \left( S + \sum_{i=1}^N \frac{s_i}{m_i} \right) n_k + \sum_{i=1}^N \frac{D_i h_i C_i}{2n_k}$$

$$C_i' \stackrel{(12)}{=} \left( S + \sum_{i=1}^N \frac{s_i}{m_i} \right) n_k + \sum_{i=1}^N \frac{D_i m_i h_i C_i}{2n_k}$$

$$\frac{dC_i'}{dn_k} = \left( S + \sum_{i=1}^N \frac{s_i}{m_i} \right) - \sum_{i=1}^N \frac{D_i m_i h_i C_i}{2n_k^2} = 0$$

$$n_k^* = \sqrt{\frac{\sum_{i=1}^N D_i h_i C_i m_i}{2 \left( S + \sum_{i=1}^N \frac{s_i}{m_i} \right)}}$$